# Zeroth-Order Online Convex Optimization

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# **Convex Optimization**

A set  $\mathcal{D}$  is convex if for every  $x, y \in \mathcal{D}$ ,  $0 \leq \lambda \leq 1$ , we have

$$\lambda x + (1 - \lambda)y \in \mathcal{D}.$$

A function  $f : D \to \mathbb{R}$  is convex if D is convex and for every  $x, y \in D$ ,  $0 \le \lambda \le 1$ ,

$$f(\lambda x + (1-\lambda)y) \leq \lambda f(x) + (1-\lambda)f(y)$$



If  $\nabla f(x^*) = 0$ , then  $x^*$  is the only global/local minima of f.

Most classical machine learning models (e.g. linear regression, logistic regression, SVM, kernel methods)

Most loss functions (e.g. mean squared loss, cross-entropy loss)

Neural networks (unfortunately they are not convex, but somehow many convex optimization algorithms still work; this is another interesting field of research)

At every timestep t, update

$$x_{t+1} = x_t - \eta g_t.$$

In gradient descent,  $g_t = \nabla f(x_t)$ .

In stochastic gradient descent,  $g_t$  is a random vector, such that  $\mathbb{E}[g_t] = \nabla f(x_t)$ , and  $\mathbb{E}[||g_t||^2] \leq G$  for some constant G.

If the constraint set  $\mathcal{D}$  is a closed and bounded convex set, we can do **projected gradient descent**, where we project  $x_{t+1}$  back to  $\mathcal{D}$  if it is "out of bound",

$$x_{t+1} = P_{\mathcal{D}}(x_t - \eta g_t).$$

### It is not hard to show that gradient descent converges. We need to assume that f is convex and $\|\nabla^2 f(x)\| < 1$

We need to assume that f is convex and  $\|\nabla^2 f(x)\| \leq L$  ("L-smooth").

## Convergence of Gradient Descent

If f is convex, then

$$f(y) \ge f(x) + \langle \nabla f(x), y - x \rangle.$$

Plugging into  $y = x^*$ ,  $x = x_t$ , we have

$$egin{aligned} f(x^*) &\geq f(x_t) + \langle 
abla f(x_t), x^* - x_t 
angle \ &= f(x_t) + rac{1}{\eta} \langle x_t - x_{t+1}, x^* - x_t 
angle. \end{aligned}$$

By "law of cosines", we get an interesting formula

$$f(x_t) \leq f(x^*) + \frac{1}{2\eta} (\|x^* - x_t\|^2 - \|x^* - x_{t+1}\|^2 + \|x_t - x_{t+1}\|^2)$$
  
$$\leq f(x^*) + \frac{1}{2\eta} (\|x^* - x_t\|^2 - \|x^* - x_{t+1}\|^2 + \eta^2 \|g_t\|^2).$$

Then, summing up from  $t = 1, \ldots, T$ , we have

$$\frac{1}{T}\sum_{t=1}^{T}f(x_t) \leq f(x^*) + \frac{1}{2\eta} \left( \|x^* - x_0\|^2 + \eta^2 \sum_{t=1}^{T} \|g_t\|^2 \right)$$

We just need to bound  $||g_t||^2!$ 

## Convergence of Gradient Descent

For gradient descent, if f is L-smooth, then by Taylor's theorem,

$$f(y) \leq f(x) + \langle 
abla f(x), y - x 
angle + rac{L}{2} \|y - x\|^2$$

Plugging into  $y = x_{t+1}$  and  $x = x_t$ , we have

$$egin{aligned} &f(x_{t+1}) \leq f(x_t) + \langle 
abla f(x_t), x_{t+1} - x_t 
angle + rac{L}{2} \|x_{t+1} - x_t\|^2 \ &\leq f(x_t) - \eta \|
abla f(x_t)\|^2 + rac{L}{2} \eta^2 \|
abla f(x_t)\|^2. \end{aligned}$$

Choosing  $\eta = \frac{1}{L}$ , we have

$$\|\nabla f(x_t)\|^2 \leq \frac{2}{\eta}(f(x_t) - f(x_{t+1}))$$

so the sum is  $O(1/\eta)$ .

The average loss is at most O(1/T).

For SGD, we just bound  $||g_t||^2$  by the variance, so sum is O(T). We can pick  $\eta = O(1/\sqrt{T})$  in this case, so the average loss is at most  $O(1/\sqrt{T})$ . Instead of a single objective function, a sequence of functions  $f_1, f_2, \ldots$  is given.

Need to pick a point  $x_t$  before knowing anything about  $f_t$ .

Try to minimize the regret compared to the minima  $x^*$  of the sum

$$R_{T} := \sum_{t=1}^{T} f_t(x_t) - \min_{x^* \in \mathcal{D}} \sum_{t=1}^{T} f_t(x^*).$$

Goal is to make  $\lim_{T\to\infty} \frac{R_T}{T} \to 0$ , so we have vanishing average regret.

(Mini-batch) stochastic gradient descent: At every epoch, sample a batch from the dataset, and optimize the loss on the batch. This is the most widely used optimization algorithm in machine learning.

**Non-convex optimization**: In empirical risk minimization, to minimize  $f(w) = \sum_{i=1}^{N} \ell(h(w, x_i), y_i)$ , with the model h non-convex and loss  $\ell$  convex, we can equivalently solve the following online convex optimization problem

$$f_t(w) = \frac{1}{N} \sum_{i=1}^N \ell(h(w_t, x_i) + \langle \nabla_w h(w_t, x_i), w - w_t \rangle, y_i).$$

**Multi-armed bandits**: There are *d* machines and you can pick which one to play. After you play, you can observe the full reward vector  $\ell_t$ . We can maintain a probability distribution  $p_t$  and define  $f_t(p_t) = -\langle \ell_t, p_t \rangle$ . Then, we can optimize using projected gradient descent.

**Reinforcement learning**: The environment changes when the player makes an action. The objective function ("value function") is dependent on the current state and the player's action.

This problem seems hard, since we don't know anything about functions in the future.

Typically, there is no guarantee that the functions are similar.

Surprisingly, gradient descent can still minimize the regret, even if we use a different function for every update step!

This is because the baseline is the minimizer of the sum  $\min_{x^* \in D} \sum_{t=1}^{T} f_t(x^*)$ , and not the minimizers of each function.

The proof is pretty much the same as the stochastic gradient case. We assume  $f_t$  to be Lipschitz, that is for every  $x, y \in D$ ,

$$|f(x)-f(y)|\leq L\|x-y\|$$

which bounds the gradient  $\|\nabla f(x)\| \leq L$ .

The regret is  $O(\sqrt{T})$ , so the average regret is vanishing (Zinkevich, 2003).

At every timestep t, instead of  $\nabla f_t(x_t)$ , we can only observe  $f_t(x_t)$ .

This is natural in scenarios where gradient is hard to obtain, or the objective function is not differentiable.

Can we still use our favorite gradient-based methods?

A company wants to decide how much money to spend on advertising their products in d channels.

The objective function is the profit, which is changing rapidly over time, and in general there is no guarantee about how it will change.

The company definitely cannot know anything about the future, and they need to choose the allocations before observing the profit.

The gradient is hard to obtain in this case — they don't know what the objective function is. They only know their profits.

We may try to approximate  $\nabla f(x)$  by

$$\langle \nabla f(x), u \rangle = \lim_{\delta \to 0} \frac{f(x + \delta u) - f(x)}{\delta}.$$

Pick a sufficiently small  $\delta$ , so when f is Lipschitz, this will be a good approximation of the partial derivatives.

In  $\mathbb{R}^d$ , we need to approximate d times to estimate the gradient.

But in the online setting, we can only observe the value of  $f_t$  at one point! Once we observe the value, the objective function changes to  $f_{t+1}$ , and we cannot get more information about  $f_t$ .

What if we can sample a unit random vector?

Key idea: Use Stoke's theorem to approximate the stochastic gradient.

$$abla f(x) pprox \mathop{\mathbb{E}}_{u \sim S} \left[ \frac{d}{\delta} f(x + \delta u) u \right] = \mathop{\mathbb{E}}_{u \sim S} \left[ \frac{d}{\delta} (f(x + \delta u) - f(x)) u \right]$$

We only need the function value at one point  $f(x + \delta v)$  to approximate the stochastic gradient!

Formally, let S be the unit sphere in  $\mathbb{R}^d$ , and  $\mathcal{B}$  be the unit ball in  $\mathbb{R}^d$ . By Stoke's theorem

$$abla \int_{\delta \mathcal{B}} f(x+v) \, dv = \int_{\delta \mathcal{S}} f(x+u) \frac{u}{\|u\|} \, du$$

Then, since  $vol(\mathcal{B}) = d vol(\mathcal{S})$ , we have

$$\mathop{\mathbb{E}}_{u\sim\mathcal{S}}\left[\frac{d}{\delta}f(x+\delta u)u\right] = \nabla\widehat{f}(x)$$

where

$$\widehat{f}(x) = \mathop{\mathbb{E}}_{v \sim \mathcal{B}}[f(x + \delta v)].$$

If we sample  $v \sim S$ , then  $\frac{d}{\delta}f(x + \delta v)v$  is a stochastic gradient of  $\hat{f}$  at x, which we can use to do stochastic gradient descent.

 $\hat{f}$  is an "average" of f in a neighborhood of radius  $\delta$ , so when  $\delta$  is small and f is Lipschitz, the difference between f and  $\hat{f}$  is small.

Algorithm: At every timestep, sample  $u_t \sim \mathcal{S}$ , and update using

$$x_{t+1} = P_{\mathcal{D}}\left(x_t - \eta \frac{\delta}{d}f_t(x_t + \delta u_t)u_t\right).$$

We can bound  $||g_t||^2 \leq \frac{d}{\delta} \max_{x \in \mathcal{D}} |f_t(x)|$ .

The standard online gradient descent proof gives us a bound on the regret of  $\hat{f}_t$  as  $O(\frac{\sqrt{T}}{\delta})$ .

However, this is a bound on the regret of  $\hat{f_t}$ ! We know that if each  $f_t$  is Lipschitz, then there is an additional error of

$$\sum_{t=1}^{T} f_t(x_t) - \sum_{t=1}^{T} \widehat{f}_t(x_t) \le O(\delta T)$$

Choosing  $\delta = O(T^{-1/4})$  gives us a regret of  $O(T^{3/4})$ .

There are various extensions of this "bandit" setting in online convex optimization.

Better algorithms exist if additional assumptions are made.

When we have two points as feedback, similarly we can estimate the gradient by

$$abla f(x) pprox \mathop{\mathbb{E}}_{u \sim \mathcal{S}} \left[ rac{d}{2\delta} (f(x + \delta u) - f(x - \delta u)) u 
ight].$$

When f is Lipschitz,  $|f(x + \delta u) - f(x - \delta u)| \le 2\delta$ , so we removed the  $\delta^{-1}$  factor in the bound of  $||g_t||!$ 

Then, we can choose  $\delta = O(\frac{1}{\sqrt{T}})$ , so regret is bounded by  $O(\sqrt{T})$ , which asymptotically is as good as when we have the gradient!

What if we sample from other than a sphere?

For a positive semidefinite matrix A, we can also approximate by

$$\nabla \widehat{f}(x) = \frac{d}{\delta} f(x + \delta A u) A^{-1} u$$

where  $u \sim S$  and

$$\widehat{f}(x) = \mathop{\mathbb{E}}_{u \sim \mathcal{B}} [f(x + \delta A u)].$$

Saha and Tewari (2011) improve the regret to  $O(T^{2/3})$  for smooth functions by finding a sequence of  $A_t$  from a self-concordant barrier function using an interior point method.

Suppose we have a repeated game, where each players have a concave utility function  $u_i$ .

The players can only optimize their action  $x_i$ , but they cannot control the action of other players  $x_{-i}$ .

We can define  $f_t(x_{t,i}) = u_i(x_{t,i}, x_{t,-i})$ , then this becomes an online optimization problem.

The individual players can use the algorithm by Flaxman et al. (2005) if they only receive their utility as feedback.

Bravo et al. (2018) shows that the game converges to a Nash equilibrium if the players use this strategy.

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